

Coding Theory

Sheet 7 Solutions

Spring 2014

1. Let C' be the set of words of even weight in the binary linear code C . Then, by Sheet 6, Exercise 5, the sum of two words of even weight also has even weight, and so C' is a linear code. Let x be any word of odd weight in C , if it exists. Then, if y is any other word of odd weight, $x + y$ has even weight and so is in C' ; that is, $y \in x + C'$. Hence

$$C = C' \cup (x + C'),$$

in which case $|C| = 2|C'|$. So, either $C' = C$ or $|C'| = \frac{1}{2}|C|$.

2. Let C be a binary $[n, k]$ code. Since

$$W_{C^\perp}(T) = 2^{-k} (1 + T)^n W_C \left(\frac{1 - T}{1 + T} \right),$$

so replacing C by C^\perp gives

$$W_C(T) = 2^{-(n-k)} (1 + T)^n W_{C^\perp} \left(\frac{1 - T}{1 + T} \right).$$

3. (a) Since

$$H = \begin{bmatrix} 1 & 0 & 1 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 \end{bmatrix},$$

so

$$G = \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix},$$

- (b) The elements of C and their weights are as follows:

$$\begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 0 & 1 & 1 & 0 & 3 \\ 0 & 1 & 0 & 0 & 1 & 1 & 3 \\ 0 & 0 & 1 & 1 & 0 & 1 & 3 \\ 1 & 1 & 0 & 1 & 0 & 1 & 4 \\ 0 & 1 & 1 & 1 & 0 & 1 & 4 \\ 1 & 0 & 1 & 0 & 1 & 1 & 4 \\ 1 & 1 & 1 & 0 & 0 & 0 & 3 \end{bmatrix}.$$

So $W_C(T) = 1 + 4T^3 + 3T^4$.

(c) Applying the MacWilliams theorem gives

$$\begin{aligned} W_C(T) &= 2^{-3} (1+T)^6 W_{C^\perp} \left(\frac{1-T}{1+T} \right) \\ &= \frac{1}{8} (1+T)^6 \left\{ 1 + 4 \left(\frac{1-T}{1+T} \right)^3 + 3 \left(\frac{1-T}{1+T} \right)^4 \right\} \\ &= \frac{1}{8} \{ (1+T)^6 + 4(1+T)^3(1-T)^3 + 3(1+T)^2(1-T)^4 \}. \end{aligned}$$

Now, this can be evaluated in various ways. Write the coefficients of the various terms:

$$\begin{array}{cccccccc} & (1+T)^6 & 1 & 6 & 15 & 20 & 15 & 6 & 1 \\ 4(1+T)^3(1-T)^3 & = & 4(1-T^2)^3 & & & & & & \\ & (1-T^2)^3 & 1 & 0 & -3 & 0 & 3 & 0 & -1 \\ & 4(1-T^2)^3 & 4 & 0 & -12 & 0 & 12 & 0 & -4 \end{array}$$

Similarly,

$$\begin{array}{cccccccc} 3(1+T)^2(1-T)^4 & 1 & -4 & 6 & -4 & 1 & & & \\ & & 2 & -8 & 12 & -8 & 2 & & \\ & & & 1 & -4 & 6 & -4 & 1 & \\ \times 3 & \hline & 1 & -2 & -1 & 4 & -1 & -2 & 1 \\ & 3 & -6 & -3 & 12 & -3 & -6 & 3 \\ & 4 & 0 & -12 & 0 & 12 & 0 & -4 \\ & 1 & 6 & 15 & 20 & 15 & 6 & 1 \\ \div 8 & \hline & 8 & 0 & 0 & 32 & 24 & 0 & 0 \\ & 1 & 0 & 0 & 4 & 3 & 0 & 0 \end{array}$$

Hence

$$W_C^\perp(T) = 1 + 4T^3 + 3T^4.$$

As a check, $W_{C^\perp}(1) = 8 = 2^3$.

(d) The elements of C^\perp and their weights are as follows:

$$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 0 & 1 & 1 & 0 & 0 & 3 \\ 1 & 1 & 0 & 0 & 1 & 0 & 3 \\ 0 & 1 & 1 & 0 & 0 & 1 & 3 \\ 0 & 1 & 1 & 1 & 1 & 0 & 4 \\ 1 & 1 & 0 & 1 & 0 & 1 & 4 \\ 1 & 0 & 1 & 0 & 1 & 1 & 4 \\ 0 & 0 & 0 & 1 & 1 & 1 & 3 \end{array}$$

So $W_C^\perp(T) = 1 + 4T^3 + 3T^4$, in agreement with the previous calculation.

(e) Thus $W_C^\perp(T) = W_C(T)$. However, $C^\perp \neq C$; but C^\perp is equivalent to C as the columns of H are a permutation of the columns of G .

4. Since C is a $[10, 7]$ code, so C^\perp is a $[10, 3]$ code. Its elements with their weights are as follows:

$$C^\perp \begin{array}{cccccccccc} 1 & 0 & 0 & 1 & 0 & 1 & 1 & 1 & 1 & 1 & 7 \\ 0 & 1 & 0 & 1 & 1 & 1 & 1 & 0 & 1 & 1 & 7 \\ 0 & 0 & 1 & 0 & 1 & 1 & 0 & 1 & 1 & 0 & 5 \\ 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 4 \\ 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 & 0 & 1 & 6 \\ 0 & 1 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 6 \\ 1 & 1 & 1 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 5 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \end{array}$$

Hence

$$W_{C^\perp}(T) = 1 + T^4 + 2T^5 + 2T^6 + 2T^7.$$

Applying the MacWilliams theorem gives

$$\begin{aligned} W_C(T) &= 2^{-3} (1+T)^{10} W_{C^\perp} \left(\frac{1-T}{1+T} \right) \\ &= \frac{1}{8} (1+T)^3 \{ (1+T)^7 + (1+T)^3 (1-T)^4 + 2(1+T)^2 (1-T)^5 \\ &\quad + 2(1+T)(1-T)^6 + 2(1-T)^7 \} \end{aligned}$$

The last four terms sum to

$$\begin{aligned} &(1-T)^4 \{ (1+T)^3 + 2(1+T)^2 (1-T) + 2(1+T)(1-T)^2 + 2(1-T)^3 \} \\ &= (1-T)^4 (7 - 3T + 5T^2 - T^3) \end{aligned}$$

Now, just writing the coefficients gives

$$\begin{array}{cccccc} 1 & -4 & 6 & -4 & 1 & & & \\ \hline 7 & -28 & 42 & -28 & 7 & & & \\ & & -3 & 12 & -18 & 12 & -3 & \\ & & & 5 & -20 & 30 & -20 & 5 \\ & & & & -1 & 4 & -6 & 4 & -1 \\ \hline 7 & -31 & 59 & -67 & 53 & -29 & 9 & -1 \end{array}$$

Putting in the term $(1+T)^7$:

$$\begin{array}{cccccc} 1 & 7 & 21 & 35 & 35 & 21 & 7 & 1 \\ \hline 8 & -24 & 80 & -32 & 88 & -8 & 16 & 0 \end{array}$$

Dividing by 8:

$$\begin{array}{cccccc} 1 & -3 & 10 & -4 & 11 & -1 & 2 & 0 \end{array}$$

Multiply this by $(1 + T)^3$:

$$\begin{array}{ccccccccccc}
 1 & -3 & 10 & -4 & 11 & -1 & 2 & 0 & 0 & 0 & 0 \\
 & 3 & -9 & 30 & -12 & 33 & -3 & 6 & 0 & 0 & \\
 & & 3 & -9 & 30 & -12 & 33 & -3 & 6 & 0 & \\
 & & & 1 & -3 & 10 & -4 & 11 & -1 & 2 & 0 \\
 \hline
 1 & 0 & 4 & 18 & 26 & 30 & 28 & 14 & 5 & 2 & 0
 \end{array}$$

Hence

$$W_C(T) = 1 + 4T^2 + 18T^3 + 26T^4 + 30T^5 + 28T^6 + 14T^7 + 5T^8 + 2T^9.$$

As a check, $W_C(1) = 128 = 2^7$.

The weight distribution of C is $(1, 0, 4, 18, 26, 30, 28, 14, 5, 2, 0)$.

5. (a) From Sheet 6, Exercise 7, all non-zero elements of C^\perp have weight $2^{r-1} = (n+1)/2$. As C^\perp is a $[2^r - 1, r]$ code, so

$$W_{C^\perp}(T) = 1 + (2^r - 1)T^{2^{r-1}} = 1 + nT^{(n+1)/2}.$$

(b)

$$\begin{aligned}
 W_C(T) &= 2^{-r} (1 + T)^n W_{C^\perp} \left(\frac{1 - T}{1 + T} \right) \\
 &= \frac{1}{n+1} (1 + T)^n \left\{ 1 + n \left(\frac{1 - T}{1 + T} \right)^{(n+1)/2} \right\} \\
 &= \frac{1}{n+1} \{ (1 + T)^n + n(1 + T)^{(n-1)/2} (1 - T)^{(n+1)/2} \}
 \end{aligned}$$

(c) $C = \text{Ham}(r, q)$ is an

$$\left[n = \frac{q^r - 1}{q - 1}, k = n - r, 3 \right]$$

code.

Let $H = [h_1, \dots, h_r]^T$ be a parity-check matrix of C with rows h_1, \dots, h_r , and let $h = \sum \lambda_i h_i$ be an element of C^\perp . If $(x_1, \dots, x_r)^T$ is the j -th column of H , then the j -th coordinate of h is zero if $\sum \lambda_i x_i = 0$. However, the number of columns (x_1, \dots, x_r) that are solutions of $\sum \lambda_i x_i = 0$ is the number N of points of $PG(r-1, q)$ in a subspace of dimension $r-2$. Hence $N = \frac{q^{r-1}-1}{q-1}$. So

$$\begin{aligned}
 w(h) = n - N &= \frac{q^r - 1}{q - 1} - \frac{q^{r-1} - 1}{q - 1} \\
 &= \frac{q^r - q^{r-1}}{q - 1} \\
 &= q^{r-1}.
 \end{aligned}$$

So

$$\overline{W}_{C^\perp}(X, Y) = X^{\frac{q^r-1}{q-1}} + (q^r - 1)X^{\frac{q^r-1}{q-1}}Y^{q^{r-1}},$$

and

$$\begin{aligned} \overline{W}_C(X, Y) &= q^{-r}\overline{W}_{C^\perp}(X + (q-1)Y, X - Y) \\ &= q^{-r}\left\{[X + (q-1)Y]^{\frac{q^r-1}{q-1}} + (q^r - 1)[X + (q-1)Y]^{\frac{q^r-1}{q-1}}(X - Y)^{q^{r-1}}\right\}. \end{aligned}$$

6. (a) The eight codewords y , $x + ty$ of C for $t \in \mathbf{F}_7$ are as follows:

1	0	4	2	3	6	x	5
0	1	4	6	5	2	y	5
1	1	1	1	1	1	$x + y$	6
1	2	5	0	6	3	$x + 2y$	5
1	3	2	6	4	5	$x + 3y$	6
1	4	6	5	2	0	$x + 4y$	5
1	5	3	4	0	2	$x + 5y$	5
1	6	0	3	5	4	$x + 6y$	5

(b) Every non-zero word in C is $\lambda x + \mu y = \lambda[x + (\mu/\lambda)y]$, when $\lambda \neq 0$, or μy when $\lambda = 0$. Hence every non-zero word in C is a multiple of one of the words in (a). So

$$A_0 = 1, \quad A_5 = 6 \times 6 = 36, \quad A_6 = 6 \times 2 = 12.$$

(c) From (b), $\overline{W}_C(X, Y) = X^6 + 36XY^5 + 12Y^6$.

(d) By the MacWilliams formula,

$$\begin{aligned} \overline{W}_{C^\perp}(X, Y) &= \frac{1}{49}\overline{W}_C(X + 6Y, X - Y) \\ &= \frac{1}{49}[(X + 6Y)^6 + 36(X + 6Y)(X - Y)^5 + 12(X - Y)^6] \\ &= X^6 + 120X^3Y^3 + 360X^2Y^4 + 972XY^5 + 948Y^6. \end{aligned}$$

Note that $\overline{W}_{C^\perp}(1, 1) = 2401 = 7^4$.

In more detail,

$$6^2 = 36, \quad 6^2 = 36, \quad 6^3 = 216, \quad 6^4 = 1296, \quad 6^5 = 7776, \quad 6^6 = 46656.$$

$(X + Y)^6$	1	6	15	20	15	6	1
$(X + 6Y)^6$	1	36	540	4320	19440	46656	46656
$(X - Y)^5$	1	-5	10	-10	5	-1	
$(X + 6Y)(X - Y)^5$	1	-5	10	-10	5	-1	0
		6	-30	60	-60	30	-6
	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>
	1	1	-20	50	-55	29	-6
$\times 36$	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>
	36	36	-720	1800	-1980	1044	-216
$12(X - Y)^6$	12	-72	180	-240	180	-72	12
$(X + 6Y)^6$	1	36	540	4320	19440	46656	46656
	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>	<hr/>
	49	0	0	5880	18140	47628	46452
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	1	0	0	120	360	972	948