

# Coding Theory

## Sheet 6

Spring 2014

- Which Hamming codes are MDS (maximum distance separable) ?
- Write out a parity-check matrix and a corresponding generator matrix for
  - Ham(2, 3);
  - Ham(2, 4);
  - Ham(3, 3);
  - Ham(3, 4);
  - Ham(3, 5);
  - Ham(4, 2).
- \* Use a parity-check matrix for Ham(4, 2), with the columns in lexicographical order, and syndrome decoding to decode
  - 00000 00000 11111;
  - 00000 11111 11111;
  - 11111 11111 11111.
- \* Let  $\mathbf{F}_4 = \{0, 1, \omega, \bar{\omega} \mid \bar{\omega} = \omega + 1 = \omega^2\}$ . Use Ham(3, 4), with a parity-check matrix having columns in lexicographical order, to decode
  - 1111111 1111111 1111111;
  - 1111111  $\omega\omega\omega\omega\omega\omega\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}\bar{\omega}$ .
- For  $x, y \in V(n, 2)$ , let
$$x \cap y = (x_1y_1, \dots, x_ny_n).$$
Show that  $w(x + y) = w(x) + w(y) - 2w(x \cap y)$ .
- If a binary  $[n, k]$  code  $C$  has parity-check matrix  $H$ , show that the extended code  $C'$  constructed in Exercise 8 of Sheet 4 has parity check matrix  $H'$ , where
$$H' = \begin{bmatrix} H & z^T \\ u & 1 \end{bmatrix},$$
with  $z = 00 \cdots 0$  of length  $n - k$  and  $u = 11 \cdots 1$  of length  $n$ .
- If  $C = \text{Ham}(r, 2)$ , show that every non-zero word of  $C^\perp$  has weight  $2^{r-1}$ .  
(Hint: Let  $H = [h_1, \dots, h_r]^T$  be a parity check matrix of  $C$  with rows  $h_1, \dots, h_r$ , and let  $h = \sum \lambda_i h_i$  be an element of  $C^\perp$ ; consider the  $j$ -th coordinate of  $h$ .)

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As part of the course assessment, hand in at the School Office solutions to the starred questions, namely 3 and 4, by 2.00 p.m. on Thursday, 20th March. Solutions to all questions will be placed online on Friday, 21st March.