

## 1.5. Logical Implication

### Definition 1.5.1. (Logical implication)

We say the logical proposition “ $r$ ” implies the logical proposition “ $s$ ” (or  $s$  logically deduced from  $r$ ) and write  $(r \Rightarrow s)$  iff  $(r \rightarrow s)$  is a tautology.

**Example 1.5.2.** Show that  $[(p \rightarrow t) \wedge (t \rightarrow q)] \Rightarrow (p \rightarrow q)$ .

**Solution.** Let P: the proposition  $(p \rightarrow t) \wedge (t \rightarrow q)$

Q: the proposition  $p \rightarrow q$

p	t	q	$p \rightarrow t$	$t \rightarrow q$	P	Q	$P \rightarrow Q$
T	T	T	T	T	T	T	T
T	T	F	T	F	F	F	T
T	F	T	F	T	F	T	T
T	F	F	F	T	F	F	T
F	T	T	T	T	T	T	T
F	T	F	T	F	F	T	T
F	F	T	T	T	T	T	T
F	F	F	T	T	T	T	T

### Remark 1.5.3.

(i) We use  $(r \Rightarrow s)$  to imply that the statement  $(r \rightarrow s)$  is true, while the statement  $(r \rightarrow s)$  alone does not imply any particular truth value. The symbol is often used in proofs as shorthand for “**implies**”.

(ii) If  $(r \Rightarrow s)$  and  $(s \Rightarrow r)$ , then we denote that by  $(r \Leftrightarrow s)$ .

**Example 1.5.4.** Show that

(i)  $(r \Rightarrow s) \equiv [(\sim r \vee s)$  is tautology].

(ii)  $(r \Leftrightarrow s) \equiv (r \equiv s)$ .

**Solution.**

(i)

1-  $(r \Rightarrow s) \equiv (r \rightarrow s)$  is tautology (Def. of  $\Rightarrow$ )

- 2-  $(r \rightarrow s) \equiv (\sim r \vee s)$  Logical Implication Law  
3-  $(\sim r \vee s)$  is tautology Inf. (1),(2)

(ii)

- 1-  $(r \Rightarrow s) \equiv (r \rightarrow s)$  is tautology }  
     $(s \Rightarrow r) \equiv (s \rightarrow r)$  is tautology } Def. of  $(\Rightarrow)$  and  $(\Leftrightarrow)$

2-  $(r \rightarrow s) \wedge (s \rightarrow r)$  is tautology.

3-  $r \leftrightarrow s$  is tautology Equivalence Law

4-  $r \equiv s$ . Inf. Remark 1.3.11

➤ Generally, the statement and its converse not necessary equivalent. Therefore,  $p \Rightarrow q$  does not mean that  $q \Rightarrow p$ .

**Example 1.5.5.** The statement “the triangle which has equal sides, has two equal legs” equivalent to the statement “the triangle which has not two equal legs has no equal sides”.

## 1.6. Quantifiers

### Definition 1.6.1.

(i) A **predicate** or **propositional function** is a statement (formula) containing variables and that may be true or false depending on the values of these variables.

- That is, a predicate is a property or relationship between objects represented symbolically.
- We represent a predicate by a letter followed by the variables enclosed between parenthesis:  $P(x)$ ,  $Q(x, y)$ , etc.

(ii) An **example** for  $P(x)$  is: value of  $x$  for which  $P(x)$  is true.

(iii) A **counterexample**  $P(x)$  is: value of  $x$  for which  $P(x)$  is false.

(iv) The set,  $X$  which contain all possible value that satisfy the formula  $P$  is called a **universal set**.

(v) A set  $Y$  which contains all values  $x$  belong to set  $X$  such that  $P(x)$  is true is called a **solution set**.

$$Y = S_p = \{x \in X: P(x) \text{ is true}\}$$

### Example 1.6.2.

(i)  $P(x) = x \leq 5 \wedge x > 3$  is true for  $x = 4$  and false for  $x = 6$  (counterexample).

(ii)  $P(x) = x \leq 5 \wedge x > 3$ , for every real numbers,  $x$  which is definitely false.

(iii) There exists an  $x$  such that  $P(x) = x \leq 5 \wedge x > 3$ , which is definitely true.

(iv) Given the statement “**Ahmad is a logician**”.

Let  $P$  represent ‘**is a logician**’ and let  $x$  represent ‘**Ahmad**’. The predicate form of this statement is  $P(x)$ . That is,  $P(x) = \text{Ahmad is a logician}$ .

(v) Let  $r$ :  $x$  is married to  $y$ .

Let  $M$  represent “**married**”. Then,  $r = M(x, y)$ .

(vi) Let  $r$ : The numbers  $x$  and  $y$  are both odd.

This statement means  $(x \text{ is odd}) \wedge (y \text{ is odd})$ .

Let  $P$  represent ‘**is a odd**’ and let  $x, y$  represent ‘**numbers**’. The predicate form of this statement is  $P(x) \wedge P(y)$ .

**Definition 1.6.3.**

(i) The phrase "for all  $x$ " ("for every  $x$ ", "for each  $x$ ") is called a **universal quantifier** and is denoted by  $\forall x$ .

(ii) The phrase "for some  $x$ " ("there exists an  $x$ ") is called an **existential quantifier** and is denoted by  $\exists x$ .

**Definition 1.6.4. (The Universal Quantifier Proposition)**

Let  $f(x)$  be a proposition function which depend only on  $x$ . A sentence  $\forall x, f(x)$  read "For all  $x, P(x)$ " mean

"For all values  $x$  in  $X$ (universal set), the predicate  $f(x)$  is true."; that is,

$$\frac{\forall x, f(x)}{\therefore f(a)}$$

**Example 1.6.5.**

(i)  $r$ : The square of all real numbers are positive.

$$r: \forall x \in \mathbb{R}, (x^2 \geq 0).$$

(ii)  $r$ : The commutative law of addition of real numbers is holed.

$$r: \forall x, \forall y \in \mathbb{R}, (x + y = y + x).$$

(iii)  $r$ : The associative law of addition of real numbers is holed.

$$r: \forall x, \forall y, \forall z \in \mathbb{R}, ((x + y) + z = x + (y + z)).$$

(iv)  $r$ : All logicians are exceptional.

Let  $L$  represent 'set of logician' and let  $E$  represent 'is exceptional'. The predicate form of this statement is  $r: \forall x \in L, E(x)$ .

(v)  $r$ : All cars are red.

Let  $X :=$  Set of cars,  $f :=$  is red. Then,  $r: \forall x \in X, f(x)$ .

### Remark .1.6.6.

(i) The "all" form, the universal quantifier, is frequently encountered in the following context:

$$\forall x (f(x) \rightarrow Q(x)),$$

which may be read,

"For all  $x$  in a universal set  $X$  satisfying  $f(x)$  also satisfy  $Q(x)$ ".

For example:

(a)  $r$ : All logicians are exceptional.

Let  $L$  represent 'is a logician' and let  $E$  represent 'is exceptional'. Then

- Predicate Logic:  $r: \forall x(L(x) \rightarrow E(x))$
- In logical English: "For all  $x$ , if  $x$  is a logician, then  $x$  is exceptional."

(b)  $r$ : The square of all real numbers are positive.

Let  $P$  represent:  $\in \mathbb{R}$  and let  $Q$  represent "is positive".

- Predicate Logic:  $r: \forall x(P(x) \rightarrow Q(x))$ ; that is,  
 $r: \forall x(\text{if } x \in \mathbb{R} \rightarrow (x^2 \geq 0)).$
- In logical English: "For all  $x$ , if  $x$  is real number, then  $x$  is positive."

(c) Every(each, any) integer number is even (or: Integer numbers are even).

Let  $P$  represent:  $\in \mathbb{Z}$  and let  $E$  represent "is even".

- Predicate Logic:  $r: \forall x(P(x) \rightarrow E(x))$ ; that is,  
 $r: \forall x(\text{if } x \in \mathbb{Z} \rightarrow E(x)).$
- In logical English: "For all  $x$ , if  $x$  is an integer, then  $x$  is even."

(ii) Parentheses are crucial here; be sure you understand the difference between the "all" form and  $\forall x, f(x) \rightarrow \forall x, Q(x)$  and  $(\forall x, f(x)) \rightarrow Q(x)$ .

### Definition 1.6.7. (The Existential Quantifier Proposition)

A sentence  $\exists x, f(x)$  read "For some  $x$ ,  $P(x)$ " or "For some  $x$  such that  $P(x)$ " mean

"For some  $x \in X$ (universal set), the predicate  $f(x)$  is true"; that is,

$$\frac{f(a)}{\therefore \exists x, f(x)}$$

**Example 1.6.8.**

(i)  $\exists x: (x \geq x^2)$  is true since  $x = 0$  is a solution. There are many others.

(ii) r: Some logicians are exceptional.

Let  $L$  represent ‘set of logician’ and let  $E$  represent ‘is exceptional’. The predicate form of this statement is  $r: \exists x \in L, E(x)$ .

(iii) r: There is a car which is red.

Let  $X :=$  Set of cars,  $f :=$  is red. Then,  $r: \exists x \in X, f(x)$ .

**Remark .1.6.9.**

(i) The “some” form, the existential quantifier, is frequently encountered in the following context:

$$\exists x (f(x) \wedge Q(x)),$$

which may be read,

“Some  $x$  in a universal set  $X$  satisfying  $f(x)$  and satisfy  $Q(x)$ ”.

For example:

(a) r: Some logicians are exceptional.

Let  $L$  represent ‘is a logician’ and let  $E$  represent ‘is exceptional’. Then

- Predicate Logic:  $r: \exists x(L(x) \wedge E(x))$
- In logical English: “For some  $x$ ,  $x$  is a logician and  $x$  is exceptional.”

(b) r: The square of some integers numbers are four (or: There is an integer for which its square is four)

Let  $P$  represent:  $\in \mathbb{Z}$  and let  $Q$  represent “is 4”.

- Predicate Logic:  $r: \exists x(P(x) \wedge Q(x))$ ; that is,

$$r: \exists x(x \in \mathbb{Z} \wedge x^2 = 4).$$

- In logical English: “For some  $x$ ,  $x$  is an integer number and  $x^2 = 4$ .”

(c) At least one integer number is even (or: Some Integers are even).

Let  $P$  represent:  $\in \mathbb{Z}$  and let  $E$  represent “is even”.

- Predicate Logic:  $r: \exists x(P(x) \wedge E(x))$ ; that is,

$$r: \exists x(x \in \mathbb{Z} \wedge E(x)).$$

- In logical English: “For some  $x$ ,  $x$  is an integer number and  $x$  is even.”

### Negation Rules of Quantifiers 1.6.10.

(i) When we negate a quantified statement, we negate all the quantifiers first, from left to right (keeping the same order), then we negative the statement.

(ii)  $\sim(x = y) = (x \neq y)$ .

(iii)  $\sim(x \equiv y) = (x \not\equiv y)$ .

(iv)  $\sim(x < y) = (y \leq x)$ .

(v)  $\sim(x \in Y) = (x \notin Y)$ .

(vi)  $\sim(\text{Even number}) = \text{Odd number}$ .

Now define the a formal universal quantifier proposition using negation.

### Definition 1.6.11.

(i)  $\forall x, f(x) = \sim \exists x, \sim f(x)$ .

(ii)  $\exists x, f(x) \equiv \sim \forall x, \sim f(x)$ .

### Example 1.6.12.

$r$ : All logicians are exceptional.

Let  $L$  represent ‘set of logician’ and let  $E$  represent ‘is exceptional’.

- Predicate Logic:  $r: \forall x \in L, E(x) = \sim \exists x, \sim E(x)$ .

- In logical English: “There is no  $x$  is a logician, for which  $x$  is not exceptional.”

### Equivalent Definitions 1.6.13.

(i)  $\sim(\forall x, f(x)) \equiv \exists x, \sim f(x)$ .

(ii)  $\sim(\exists x, f(x)) \equiv \forall x, \sim f(x)$ .

- (iii)  $\sim [\forall x (f(x) \rightarrow Q(x))] \equiv \exists x (f(x) \wedge \sim Q(x))$   
 $\equiv$  Some  $f(x)$  are not  $Q(x)$
- (iv)  $\sim (\exists x, (f(x) \wedge Q(x))) \equiv \forall x, \sim f(x) \vee \sim Q(x) \equiv \forall x (f(x) \rightarrow \sim Q(x))$   
 $\equiv$  No  $f(x)$  are  $Q(x)$

### Example 1.6.14.

(i) Express each of the following sentences in the form  $\forall x, P(x)$  and then give its negation in both cases  $\forall x, P(x)$  and in words.

r: The square of every real number is non-negative.

**Solution.**

- **$\forall x, P(x)$  form:** r:  $\forall x \in \mathbb{R}, x^2 \geq 0$ .
- **Negation:**  $\sim r: \sim (\forall x \in \mathbb{R}, x^2 \geq 0) \equiv \exists x \in \mathbb{R}, \sim (x^2 \geq 0) \equiv \exists x \in \mathbb{R}, x^2 < 0$ .
- **Negation in words:**  $\sim r$ : There exists a real number whose square is negative.

(ii) Let **r: Student who is intelligent will succeed**. Write “ r ” in predicate logic and English logic, and then give its negation in both cases.

**Solution.**

Let P: Student

Q: intelligent.

S: Succeed

- **Predicate Logic:** r:  $\forall x ((P(x) \wedge Q(x)) \rightarrow S(x))$
- **Negation:**  $\sim r: \sim [\forall x ((P(x) \wedge Q(x)) \rightarrow S(x))]$   
 $\equiv \sim [\forall x (\sim (P(x) \wedge Q(x)) \vee S(x))]$  Implication Law.  
 $\equiv \exists x ((P(x) \wedge Q(x)) \wedge \sim S(x))$  De Mover's Law.
- **English logic:**  $\sim r$ : There exist student who is intelligent and not succeed.

(iii) r: Some integer numbers are even but not odd.

Let  $\mathbb{Z} :=$  Set of Integers,  $f :=$  is even,  $P :=$  is odd.

- **Predicate Logic:** r:  $\exists x \in \mathbb{Z}, (f(x) \wedge \sim P(x)) \equiv \sim [\forall x (f(x) \rightarrow P(x))]$ .
- **English Logic:** r: Not all even integers are odd.
- **Negation:**  $\sim r: \sim \sim [\forall x (f(x) \rightarrow P(x))] = [\forall x (f(x) \rightarrow P(x))]$ .
- **Negation in words:** All even integer numbers are odd.



### Remark 1.6.15.

Sometimes the English sentences are **unclear** with respect to quantification, or in another words, quantified statements are often misused in **casual (informal) conversation**.

For example:

(i) “If you can solve any problem we come up with, then you get an A for the course”

The phrase “you can solve any problem we can come up with” could reasonably be interpreted as either a universal or existential quantification:

(a) “you can solve every problem we come up with”,

(b) “you can solve at least one problem we come up with”.

(ii) r: All students do not pay full tuition.

Here “ r ” could reasonably be interpreted as

(a) Not all students pay full tuition (Or: There exist some students do not pay full tuition).

(b) No students are pay full tuition (Or: There are no students are pay full tuition).

**Mathematical context:** Not all students pay full tuition.

(iii) r: All integer numbers are not even.”

(a) Not all integer numbers are even.

(b) No integer numbers are even (Or: There are no even integers).

**Mathematical context:** Not all integer numbers are even.

**Combined Quantifiers 1.6.16.** There are six ways in which the quantifiers can be combined when two variables are present:

(1)  $\forall x \forall y, f(x, y) \equiv \forall y \forall x, f(x, y)$  = For every  $x$ , for every  $y$ ,  $f(x, y)$ .

(2)  $\forall x \exists y, f(x, y) \equiv$  For every  $x$ , there exists a  $y$  such that  $f(x, y)$ .

(3)  $\forall y \exists x, f(x, y) \equiv$  For every  $y$ , there exists an  $x$  such that  $f(x, y)$ .

(4)  $\exists x \forall y, f(x, y) \equiv$  There exists an  $x$  such that for every  $y$ ,  $f(x, y)$ .

(5)  $\exists y \forall x f(x, y) \equiv$  There exists a  $y$  such that for every  $x$ ,  $f(x, y)$ .

(6)  $\exists x \exists y, f(x, y) \equiv \exists y \exists x, f(x, y) =$  There exists an  $x$  such that there exists a  $y$ ,  $f(x, y)$ .

### Example 1.6.17.

(i)  $r: \exists x \in \mathbb{R} \exists y \in \mathbb{R} : P(x, y) = (x^2 + y^2 = 2xy)$ . The proposition “ $r$ ” is true since  $x = y = 1$  is one of many solutions.

(ii)  $s: \forall x \in \mathbb{R} \exists y \in \mathbb{R} : P(x, y) = (y^3 = x)$ . The proposition “ $s$ ” is true since  $y = \sqrt[3]{x}$  is solution for  $P(x, y)$ .

(iii)  $s: \exists x \in \mathbb{R} \forall y \in \mathbb{R} : P(x, y) = (y^3 = x)$ . Here “ $s$ ” mean there is an “ $x$ ” real such that for every “ $y$ ” real,  $P(x, y)$  is true. The proposition “ $s$ ” is not true since no real numbers have this property.

(iv)  $r$ : For all  $x$ , there exists  $y$  such that  $xy = 1$ .

#### Solution.

- $\forall x, P(x)$  form:  $r: \forall x, \exists y$  such that  $xy = 1$ .
- **Negation:**  $\sim r: \sim (\forall x, \exists y \text{ such that } xy = 1)$   
 $\equiv \exists x, \sim (\exists y \text{ such that } (xy = 1))$   
 $\equiv \exists x, \forall y \text{ such that } xy \neq 1$ .
- **Negation in words:**  $\sim r$ : There exists  $x$  such that for all  $y$ ,  $xy \neq 1$ .

(v) The following are equivalents.

(a)  $\sim[\forall x \forall y, f(x, y)] \equiv \exists x \exists y, \sim f(x, y)$ .

(b)  $\sim[\exists x \exists y, f(x, y)] \equiv \forall x \forall y, \sim f(x, y)$ .

(c)  $\sim[\forall x \exists y, f(x, y)] \equiv \exists x \forall y, \sim f(x, y)$ .

(d)  $\sim[\exists x \forall y, f(x, y)] \equiv \forall x \exists y, \sim f(x, y)$ .

#### Solution. Exercise.

## 1.7. Logical Reasoning

### Definition 1.7.1. (Arguments)

An **argument** is a series of statements starting from hypothesis (premises/assumptions) and ending with the conclusion.

From the definition, an argument might be valid or invalid.

### Definition 1.7.2. (Valid Arguments)(Proofs)

An argument is said to be **valid** if the hypothesis implies the conclusion; that is, if  $s$  is a statement implies from the statements  $s_1, s_2, \dots, s_n$ , then write as

$$s_1, s_2, \dots, s_n \mapsto s.$$

**Note 1.7.3.** In mathematics, the word **proof** is used to mean simply a valid argument. Many proofs involve more than two premises and a conclusion.

### Example 1.7.4.

(i) Let  $s_1$ : Some mathematicians are engineering

$s_2$ : Ali is mathematician

$s$ : Ali is engineering

Show that the argument is valid.

#### Solution.

The argument  $s_1, s_2 \mapsto s$  is not valid, since not all mathematicians are engineering.

(ii) Let  $s_1$ : There is no lazy student

$s_2$ : Ali is artist

$s_3$ : All artist are lazy

Find a conclusion  $s$  for the above premises making the argument  $s_1, s_2, s_3 \mapsto s$  is valid.

#### Solution.

Ali is-----.

### Remark 1.7.5.

(i) An argument

is valid if and only if

$$s_1, s_2, \dots, s_n \mapsto s$$

$(s_1 \wedge s_2 \wedge \dots \wedge s_n) \rightarrow s$  is tautology; that is,

$$(s_1 \wedge s_2 \wedge \dots \wedge s_n) \Rightarrow s.$$

(ii) An argument does not depend on the truth of the premises or the conclusion but it just interested only in the question

**“Is the conclusion implied by the conjunction of the premises?”**