

Figure 3.5 Critical Path ( $j_3-j_4-j_5-j_8-j_7$ ) for the precedence diagram.

### MINIMIZATION OF MAKESPAN WITH $P_j = 1$ AND PRECEDENCE PROBLEM

( $P_m / P_j = 1, tree / C_{max}$ )

This parallel machine-scheduling problem pertains to jobs having precedence relationship represented by a *tree*. All the jobs on the tree have process time equal to unity. If the jobs have precedence relationship described either by *intree* or by *outtree* (shown below), then CP rule can be applied to minimize  $C_{max}$ .

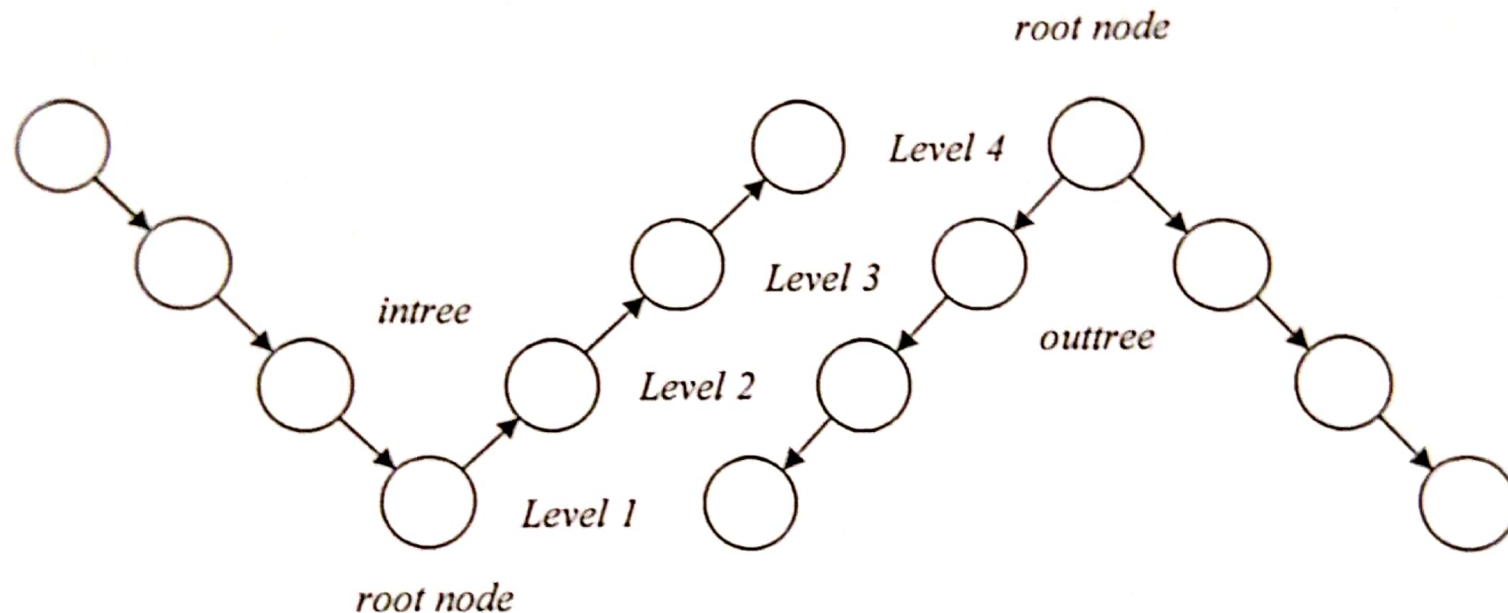


Figure 3.6 Relationship of *intree* and *outtree*

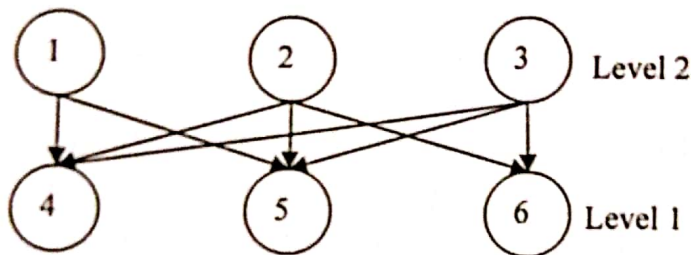
#### 3.4.1 CP Rule

The root node for *intree* precedence graph is at lowest level (Level 1), whereas root node for *outtree* precedence graph is at highest level (Level 5). CP rule imparts high priority to a job having longest string of jobs in the precedence graph. This means that a job having highest level in the precedence graph will be scheduled first. Hence CP rule is in fact **Highest Level First** rule. This means that root node of *outtree* precedence graph will have highest priority and, root node of *intree* precedence graph will be given least priority. Note that upper bound (UB) on the worst case performance of CP rule is given by the following formula.

$$\frac{C_{\max}(\text{CP})}{C_{\max}(\text{OPT})} \leq \frac{4}{3}$$

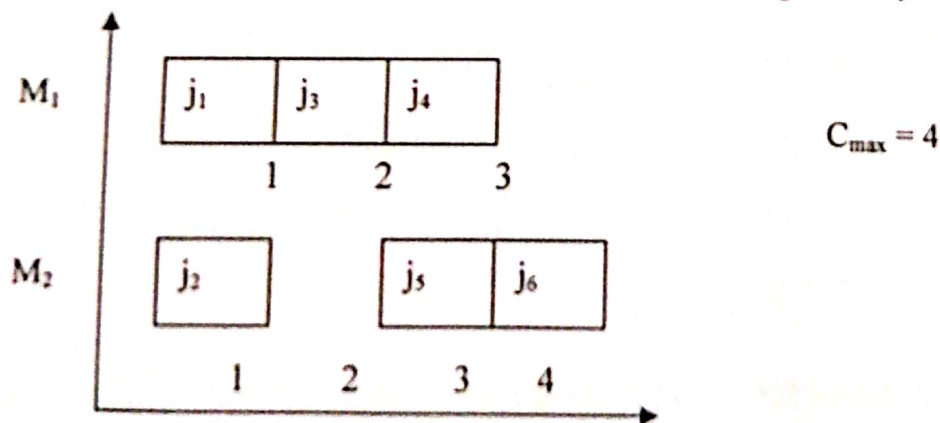
**Example 3.4**

For the  $Pm | p_j = 1, \text{tree} | C_{\max}$ , precedence graph shown below, apply CP rule and find  $C_{\max}$ . The production shop has two machines in parallel.



**Solution:**

Using CP rule, first select jobs with Highest Level, hence the jobs in candidate set at time  $t=0$  are from Level 2 jobs; i.e., jobs  $\{j_1, j_2, j_3\}$ . Schedule job 1 on machine 1 and job 2 on machine 2. At time 1, only job 3 is schedulable. Hence, schedule job 3 on machine 1 at time 1. At time 2, all the jobs at level 1 are schedulable. Jobs at Level 1 include  $\{j_4, j_5, j_6\}$ . However, only two jobs can be scheduled. So, schedule job 4 and 5 on machine 1 and 2 respectively at time 2. Finally, at time 3 schedule job 6 on machine 2. The complete schedule is shown in Gantt chart (Figure 3.7).



**Figure 3.7** Gantt chart for schedule generated by CP Rule.

### 3.4.2 LNS Rule

The LNS rule is stated as follows:

“Find the number of successors (NS) of all jobs in the precedence graph. While scheduling jobs on parallel machines, give priority to jobs having largest number of successors.”

#### Example 3.5

Solve the problem in Example 3.4 using LNS rule.

#### Solution

The number of successors for each job in example 3.5 according to precedence graph is shown in the following table:

Job (j)	1	2	3	4	5	6
$P_j$	1	1	1	1	1	1
NS	2	3	3	0	0	0

At time  $t=0$ ,

Assign jobs  $j_2$  and  $j_3$  to machines  $M_1$  and  $M_2$ .

At time  $t=1$ ,

Schedule job  $j_1$  on machine  $M_1$ , machine  $M_2$  is also free and ready at time  $t=1$ .

All the three jobs  $\{j_3, j_4, j_5\}$  have  $NS=0$ . But only job  $j_6$  is schedulable at  $t=1$  because it satisfies the precedence constraint. Hence, schedule job  $j_6$  at time  $t=1$ .

At time  $t=2$ ,

The remaining unscheduled jobs  $\{j_4, j_5\}$  are schedulable. So, schedule jobs  $j_4$  and  $j_5$  on machines  $M_1$  and  $M_2$  respectively at time  $t=2$ . Note, the makespan ( $C_{max}$ ) of this schedule is 3 time units.

The complete schedule is presented in Gantt chart (Figure 3.8)

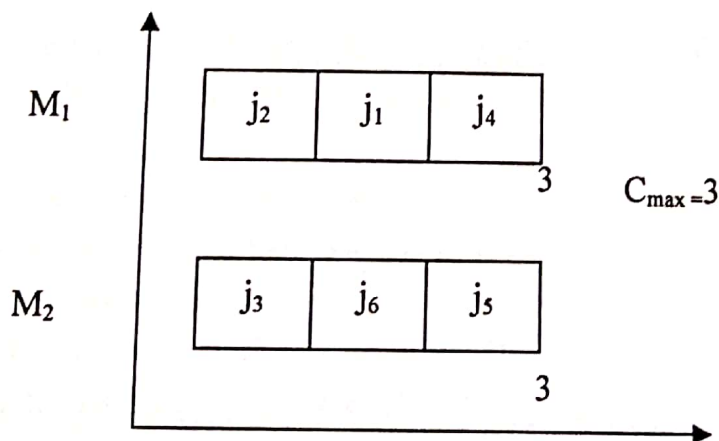


Figure 3.8 Gantt chart for Optimal Schedule generated by LNS rule.