## CHAPTER FOUR

## CRYPTANALYSIS OF TRANSPOSITION

## CIPHER PROBLEMS USING COMBINATORIAL

## OPTIMIZATION PROBLEMS TECHNIQUES

### 4.1 Terminology

- Cryptography: is the study of principles and techniques by which information can be concealed in ciphertexts and later revealed by legitimates users employing the secret key. Its concern Encryption and Decryption processes
- Cryptanalysis: is the science (and art) of recovering information from ciphertexts without knowledge of the key.
- Encryption: is a process of encoding a message so that the meaning of the message is not obvious.
- Decryption: is the reverse process: transforming an encrypted message back into its normal form.
- Cryptosystem: A system for encryption and decryption.
- The original form of a message is known as Plaintext, and the encrypted form is called Ciphertext.


### 4.2 Notations

- $M$ : plaintext message, $P=\left[m_{l}, m_{2}, \ldots, m_{n}\right]$.
- $C$ : ciphertext can be written as $C=\left[c_{1}, c_{2}, \ldots, c_{m}\right]$.
- $E$ : is the encryption algorithm.
- $D$ is the decryption algorithm.
- the transformations between P and C are $C=E(M)$ and $M=D(C)$, so $M=D(E(M))$.
- K: key, so that the $C=E(K, M)$. and $M=D(K, E(K, M))$.


### 4.3 Simple Transpositions

The goal of transposition is diffusion, spreading the information from the message or the K out widely across the C . Because a transposition is a rearrangement of the symbols of a message, it is also known as a permutation.

The columnar transposition is a rearrangement of the characters of the plaintext into columns.

The following example is a five-column transposition. The plaintext characters are separated into blocks of five and arranged one block after another, as shown here.

| $\mathrm{c}_{1}$ | $\mathrm{c}_{2}$ | $\mathrm{c}_{3}$ | $\mathrm{c}_{4}$ | $\mathrm{c}_{5}$ |
| :--- | :--- | :--- | :--- | :--- |
| $\mathrm{c}_{6}$ | $\mathrm{c}_{7}$ | $\mathrm{c}_{8}$ | $\mathrm{c}_{9}$ | $\mathrm{c}_{10}$ |
| $\mathrm{c}_{11}$ | $\mathrm{c}_{12}$ | etc. |  |  |

The resulting C is formed by transversing the columns.
$\mathrm{c}_{1} \mathrm{c}_{6} \mathrm{c}_{11} \ldots . \mathrm{c}_{2} \mathrm{c}_{7} \mathrm{c}_{12} \ldots . . \mathrm{c}_{3} \mathrm{c}_{8}$, etc.
Example (4.1): you would write the plaintext message as:

| T | H | I | S | I |
| :---: | :---: | :---: | :---: | :---: |
| S | A | M | E | S |
| S | A | G | E | T |
| O | S | H | O | W |
| H | O | W | A | C |
| O | L | U | M | N |
| A | R | T | R | A |
| N | S | P | O | S |
| I | T | I | O | N |
| W | O | R | K | S |

The resulting ciphertext would then be read off as:
tssoh oaniw haaso lrsto imghw utpir seeoa mrook istwc nasns

The length of this message happened to be a multiple of five, so all columns came out the same length.

Let E and $\mathrm{D}=\mathrm{E}^{-1}$ be encryption and decryption function of TCP respectively. The ciphertext $C_{m}$ of $T C P$, where $1 \leq m \leq n$ !, using arbitrary encryption key $\mathrm{EK}_{\mathrm{m}}$ with length n is:

$$
\begin{equation*}
\mathrm{C}_{\mathrm{m}}=\mathrm{E}\left(\mathrm{M}, \mathrm{EK}_{\mathrm{m}}\right) \tag{E}
\end{equation*}
$$

Let $\mathrm{DK}_{\mathrm{m}}$ be the decryption key corresponding to the $\mathrm{EK}_{\mathrm{m}}$ ( $\sigma$ of nsequence) for ciphertext $\mathrm{C}_{\mathrm{m}}$ of TCP and $\mathrm{P}_{\mathrm{m}}$ be the decrypted text using $\mathrm{DK}_{\mathrm{m}}$, is:

$$
\begin{equation*}
\mathrm{M}=\mathrm{M}_{\mathrm{m}}=\mathrm{D}\left(\mathrm{C}_{\mathrm{m}}, \mathrm{DK}_{\mathrm{m}}\right) \tag{D}
\end{equation*}
$$

Its clear that $C_{m}$ (and $M_{m}$ ) consists of $n$ columns.

Example (4.2): Let's have the following PT message (showed in uppercase letters):

| 1 | 2 | 3 | 4 |
| :---: | :---: | :---: | :---: |
| T | H | E | Q |
| U | I | C | K |
| B | R | O | W |
| N | F | O | X |
| J | U | M | P |
| S | O | V | E |
| R | T | H | E |
| L | A | Z | Y |
| D | O | G | X |

The size of the permutation is known as the period. For this example a simple transposition cipher with a period of 4 is used. Let $\Pi=(3,1,4,2)$ be encryption key. Then the message is broken into blocks of 4 characters. Upon encryption the $3^{\text {rd }}$ character in the block will be moved
to position 1 , the $1^{\text {st }}$ to position 2 , the $4^{\text {th }}$ to position 3 and the $2^{\text {nd }}$ to position 4.

| 3 | 1 | 4 | 2 |
| :---: | :---: | :---: | :---: |
| e | t | q | h |
| c | u | k | i |
| o | b | w | r |
| o | n | x | f |
| m | j | p | u |
| v | s | e | o |
| h | r | e | t |
| z | l | y | a |
| g | d | x | o |

The resulting ciphertext (in lowercase letters) would then be read off as:
etqhc ukiob wronx fmjpu vseoh retzl yagdx o
Notice also that decryption can be achieved by following the same process as encryption using the "inverse" of the encryption permutation. In this case the decryption key (DK), $\Pi^{-1}$ is equal to $(2,4,1,3)$.

