

Lecture Three

Regular Expression (RE)

Regular Expression is a set of symbols, thus if alphabet = {a, b}, then aab, a, baba, bbbbbb, and baaaaa would all be strings of symbols of alphabet.

In addition, we include an empty string denoted by Λ which has no symbols in it. We now introduce the use of the Kleene star applied not to a set but directly to the letter x and written as (x^*) . The simple expression x^* will be used to indicate some sequence of x's.

$$x^* = \{ \Lambda, x, x^2, x^3, x^4, \dots \} = \{ x^n \text{ for } n = 0, 1, 2, 3, 4, \dots \}$$

Examples:

- 1- $(ab)^* = \{ \Lambda, ab, abab, ababab, \dots \}$
- 2- $ab^*a = \{ aa, aba, abba, abbba, abbbbba, \dots \}$
- 3- $a^*b^* = \{ \Lambda, a, b, aa, ab, bb, aaa, aab, abb, bbb, aaaa, \dots \}$

Notice that **ba** and **aba** are not in this language. Also we should be very careful to observe that $a^*b^* \neq (ab)^*$

$$4- L_1 = \{ x^{\text{odd}} \} = x(xx)^* \text{ or } (xx)^*x = \{ x, xxx, xxxxx, \dots \}$$

$$5- L_2 = \{ x^{\text{even}} \} = xx(xx)^* \text{ or } (xx)^*xx \text{ or } (xx)^* = \{ \Lambda, xx, xxxx, \dots \}$$

- 6- Consider the language L_3 defined over the alphabet $\Sigma = \{a, b, c\}$, All the words in L_3 begin with an **a** or **c** and then are followed by some number of **b**'s. We may write this as:

$$L_3 = (a + c)b^*$$

- 7- Consider a finite language L_4 that contains all the strings of **a**'s and **b**'s of length exactly three.

$$L_4 = \{ aaa, aab, aba, abb, baa, bab, bba, bbb \}$$

So we may write:

$$L_4 = (a + b)(a + b)(a + b) \text{ or } (a + b)^3$$

In general, if we want to refer to the set of all possible strings of **a**'s and **b**'s of any length, we could write:

$$L_4 = (a + b)^*$$

- 8- Construct RE for all words that begin with the letter **a**: $a(a + b)^*$
- 9- All words that begin with an **a** and end with **b** can be defined by the expression: $a(a + b)^*b$
- 10- The language of all words that have at least two **a**'s can be described by the expression:

$$(a + b)^*a(a + b)^*a(a + b)^* \quad \underline{\text{or}} \quad b^*ab^*a(a + b)^*$$

- 11- The language of all words that have at least one **a** and at least one **b**:

$$(a + b)^*a(a + b)^*b(a + b)^* \quad \underline{\text{or}} \quad bb^*aa^*$$

- 12- The words of the form some **b**'s followed by some **a**'s. These exceptions are all defined by the regular expression: $bb^*aa^* \equiv b^+a^+$ Homework:

- 1- Find a regular expression over the alphabet {a, b}:

- a. $L_1 = \{\text{All strings that contain exactly three a's}\}$
- b. $L_2 = \{\text{All strings that end with ab}\}$
- c. $L_3 = \{\text{All strings in which letter a is even number}\}$
- d. $L_4 = \{\text{All strings that contain exactly two successive a's.}\}$

- 2- Find the output (words) for the following regular expressions:

- a. aa^*b
- b. bba^*a
- c. $(a + b)^*ba$
- d. $(0+1)^*00(0+1)^*$
- e. $(11 + 0)^*(0+11)^*$
- f. $01^* + (00+101)^*$
- g. $(a+b)^*abb^+$
- h. $((01+10)^*11)^*00)^*$